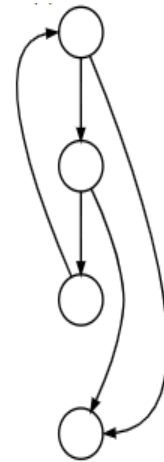


# Program Analysis

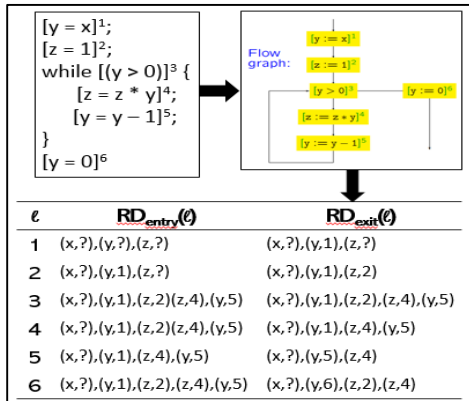
Venkatesh Vinayakarao

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Mar – Apr, 2018  
Chennai Mathematical Institute

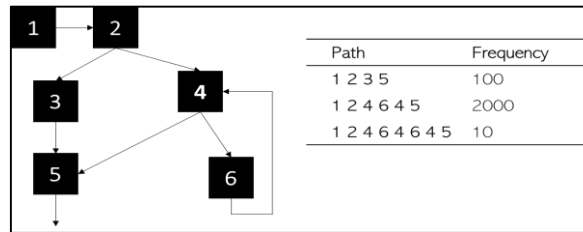


*“Popularity of data-driven software engineering has led to an increasing demand on the infrastructures to support efficient execution of tasks that require **deeper source code analysis**.”* – Ganesha and Hridesh, ICSE 2018.

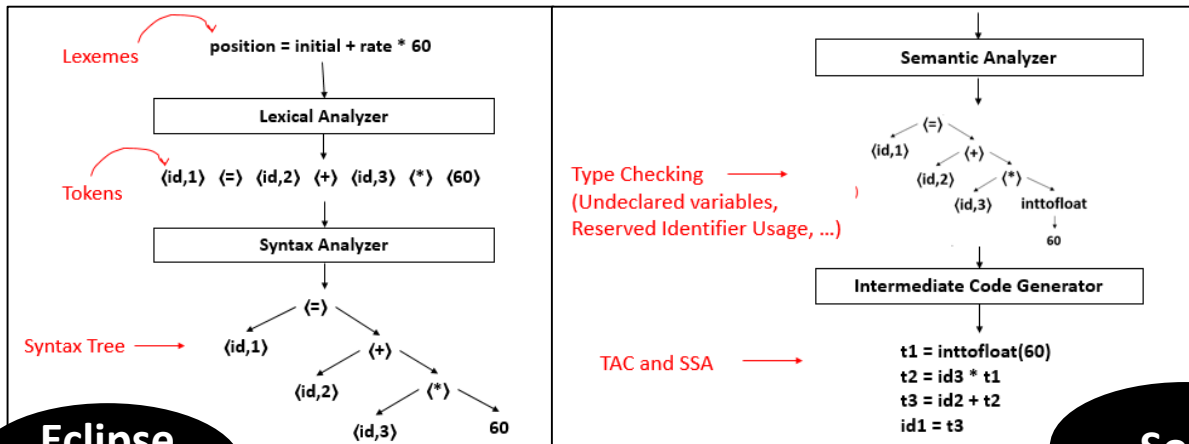
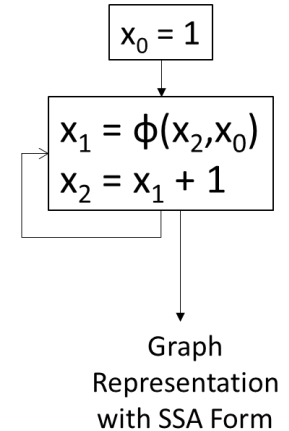
# Quick Review



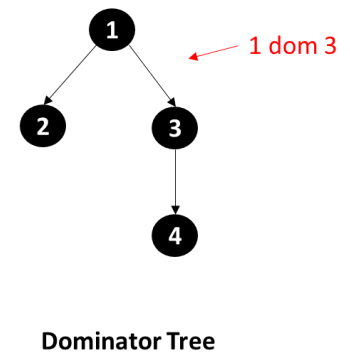
**Static Analysis**



**Dynamic Analysis**



**Structure of a Compiler**



**Derived Program Representations**

# Agenda

- Program Representations
  - Basic Blocks
  - Control Flow Graphs
  - Static Single Assignment
  - Three-Address-Code
- Hands-On Session on Traversing a CFG
- More Program Representations
  - Data Dependence
  - Control Dependence
- Dynamic Views of a Program

# Assignment-1

---

Deadline: Submit your code on Moodle by 22<sup>nd</sup> March 2019 21:00 Hrs. Include a ReadMe.txt file with instructions on how to execute your code.

# Data Dependence

S1 **x = y + 10;**

S2 **z = x \* y;**

The execution order of S1  
and S2 cannot be  
changed. S2 is data  
dependent on S1.

# Control Dependence Graph

**y is control-dependent on x**  
x determines whether y will be executed

```
S1  if (x > 0) {  
S2      y = 10;  
S3  }
```

S1 determines whether S2 will be executed. S2 is control dependent on S1.

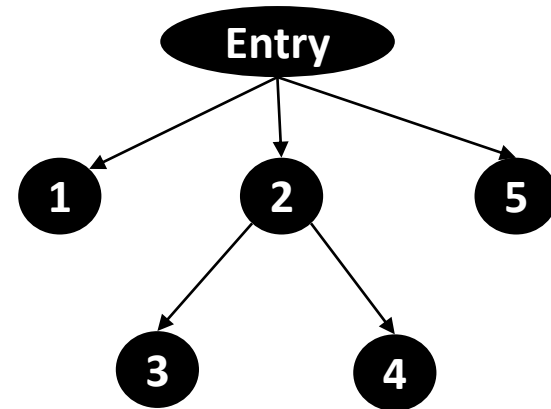
```
S1  while (y > 10) {  
S2      y = y + 1;  
S3  }
```

S1 determines whether S2 will be executed.  
S2 is control dependent on S1.  
Also, S1 is control dependent on itself!

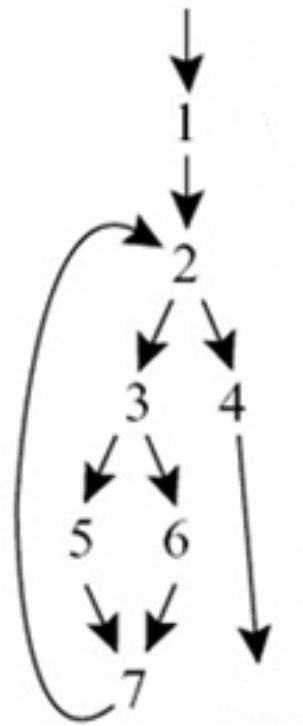
# Control Dependence Graph

```
x = 10;  
if (x > k)  
    y = 10;  
else  
    y = 1;  
print y;
```

```
[x = 10]1;  
if [(x > k)]2  
    [y = 10]3;  
else  
    [y = 1]4;  
[print y]5;
```



# Graded Quiz: Draw the Control Dependence Graph for this CFG

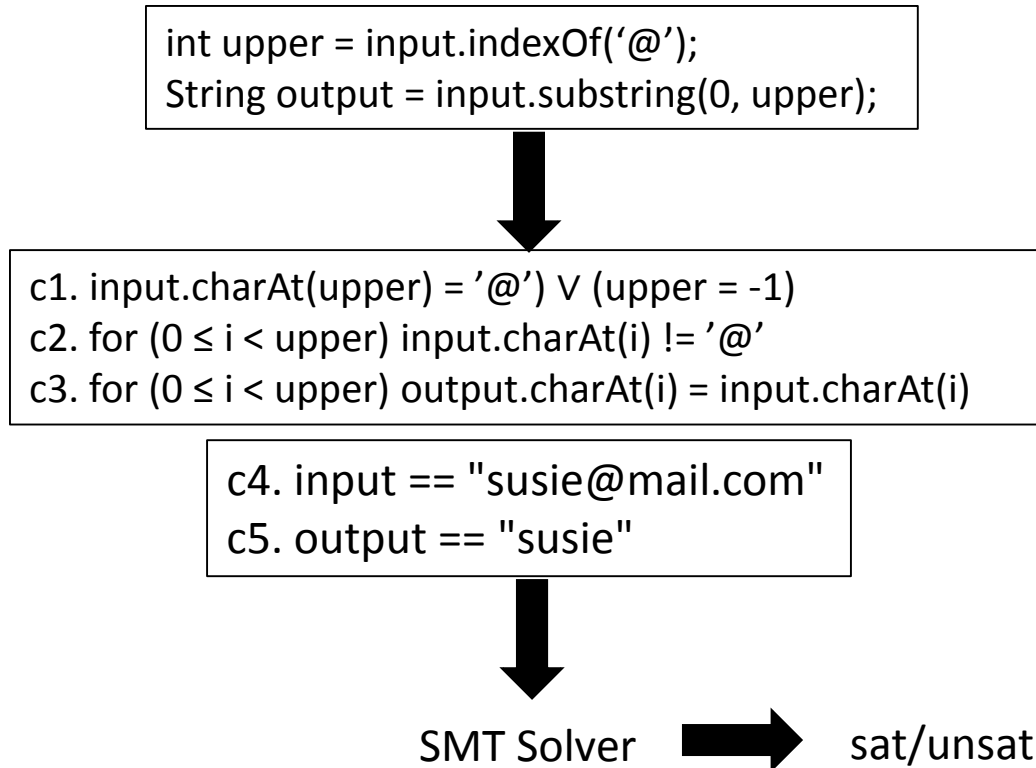


Control Flow Graph





# Program as Constraints



# SMT Solvers

- Satisfiability Modulo Theory
  - Input
    - a formula over a background theory
  - Output
    - sat/unsat
      - If sat, it can also compute a satisfying assignment
- We use z3 (SMT) Solver
  - Several theories supported
    - QF\_LIA: quantifier free linear integer arithmetic
    - QF\_IDL: quantifier free integer difference logic
    - ...

# Z3 from Java

```
Model check(Context ctx, BoolExpr f, Status sat) throws TestFailedException
{
    Solver s = ctx.mkSolver();
    s.add(f);
    if (s.check() != sat)
        throw new TestFailedException();
    if (sat == Status.SATISFIABLE)
        return s.getModel();
    else
        return null;
}
```

---

<https://github.com/Z3Prover/z3/blob/master/examples/java/JavaExample.java>

# Readings

- Please read
  - Dragon Book
    - Section 6.2 for Three-Address Code and Static Single-Assignment.
    - Section 8.4.1 and 8.4.3 for Basic Blocks and Flow Graphs.
    - Section 9.6.1 for Dominators and Dominator Tree.
  - Program as Constraints
    - [Solving the Search for Source Code](#), Kathryn Stolee (Optional)

# Runtime View of a Program

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# JProfiler Demo

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<https://www.ej-technologies.com/blog/2017/03/finding-a-memory-leak-with-jprofiler/>

<https://www.youtube.com/watch?v=032aTGa-1XM>



Telemtries



Live memory



Heap walker



CPU views



Threads



Monitors &amp; locks



Databases




JEE &amp; Probes




MBeans

JProfiler

**No snapshot has been taken.****For a maximum of features:**


Press  to take a JProfiler heap snapshot

- The snapshot is displayed in this frame and saved together with profiling information from other views
- For live profiling sessions, special features are available
- Integrations with other views require this snapshot type

Press  to indicate the starting point of a use case

- All objects that are currently on the heap will be marked as old
- When you take the next heap snapshot, new and old objects will be listed separately in the header
- You can select new or old objects only, making it easy to track down memory leaks

**For a minimum of overhead:**

Press  to take an HPROF heap snapshot

- The snapshot is saved separately and displayed in another frame
- Not all features are available
- Memory and CPU overhead in the profiled VM are lower than for the JProfiler snapshot



0 active recordings

VM #1

Profiling

# Heap Snapshot (using JProfiler)

Animated Bezier Curve Demo - JProfiler 10.15

Session View Profiling Window Help

Start Center Stop Save Snapshot Session Settings Start Recordings Stop Recordings Start Tracking Run GC Add Bookmark Export View Settings Help Take Snapshot Mark Heap Back Forward Go To Start Show Selection

Telemetries Live memory Heap Walker CPU views Threads Monitors & locks Databases JEE & Probes MBeans

Classes Allocations Biggest Objects References Time Inspections Graph

Current object set: 72,887 objects in 1,259 classes  
1 selection step, 5,096 kB shallow size

Classes Use ... Group By Class Loaders Calculate estimated retained sizes

Name	Instance Count	Size
byte[]	13,030	665 kB
java.lang.String	11,562	277 kB
java.util.HashMap\$Node	5,955	190 kB
java.lang.Object[]	5,575	284 kB
java.lang.invoke.LambdaForm\$Name	3,173	101 kB
java.lang.Class	3,020	966 kB
java.util.concurrent.ConcurrentHashMap\$Node	2,659	85,088 bytes
float[]	1,263	221 kB
java.lang.Long	1,252	30,048 bytes
java.awt.geom.GeneralPath	1,246	39,872 bytes
java.lang.invoke.MemberName	1,136	54,528 bytes
java.util.Hashtable\$Entry	1,059	33,888 bytes
java.lang.Integer	800	12,800 bytes
java.lang.invoke.ResolvedMethodName	759	18,216 bytes
java.util.HashMap	658	31,584 bytes
java.lang.Class[]	625	23,064 bytes
java.util.HashMap\$Node[]	592	77,968 bytes
<b>Total:</b>	<b>72,887</b>	<b>5,096 kB</b>

Class View Filters

Selection step 1: All objects after full GC, retaining soft references  
72,887 objects in 1,259 classes



# Call Tree

The screenshot displays the IntelliJ IDEA memory profiler interface. At the top is a toolbar with icons for Session, Profiling, and View specific actions. Below the toolbar is a sidebar with navigation options: Telemetries, Live memory, All Objects, Recorded Objects, Allocation Call Tree (selected), Allocation Hot Spots, and Class Tracker. The main area shows the Allocation Call Tree for 'Live objects at 15:21, All classes' with an aggregation level of 'Methods'. The tree lists various methods with their respective allocation percentages, sizes, and counts.

Recorded allocations of: Live objects at 15:21, All classes

Aggregation level: **m** Methods

Method	Allocation %	Allocation Size	Allocation Count
java.awt.EventQueue.run	93.2%	2,024 kB	33,310 alloc.
bezier.BezierAnim\$Demo.paint	55.1%	1,197 kB	14,938 alloc.
bezier.BezierAnim\$Demo.drawDemo	46.8%	1,017 kB	12,533 alloc.
java.awt.Graphics2D.fill	40.6%	882 kB	10,609 alloc.
java.awt.geom.GeneralPath.<init>	5.5%	119 kB	1,443 alloc.
java.awt.Graphics2D.draw	0.7%	15,392 bytes	481 alloc.
bezier.BezierAnim\$Demo.createGraphics2D	7.8%	169 kB	1,924 alloc.
bezier.BezierAnim\$Demo.run	6.8%	148 kB	4,411 alloc.
bezier.BezierAnim\$Demo.scheduleBlockingActivity	3.5%	76,880 bytes	2,402 alloc.
java.awt.EventQueue.invokeLater	3.2%	69,184 bytes	1,921 alloc.
bezier.BezierAnim\$Demo\$1.<init>	0.4%	7,696 bytes	481 alloc.
bezier.BezierAnim\$Demo.scheduleRepaint	3.3%	71,984 bytes	2,009 alloc.

# Summary

- Program Representations
  - Lexemes
  - Tokens
  - Abstract Syntax Tree
  - Control Flow Graphs
  - Three-Address Code
  - Static Single Assignment
  - Dominator Tree
  - Data Dependence Graph
  - Control Dependence Graph
  - Program as Constraints
  - Runtime Representations
  - Program Dependence Graph, Hammocks Graphs, State Machines, Petri Nets ...