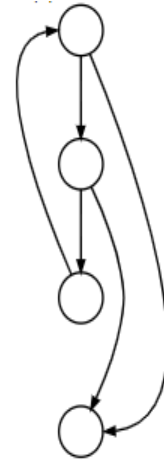


Program Analysis

Venkatesh Vinayakarao

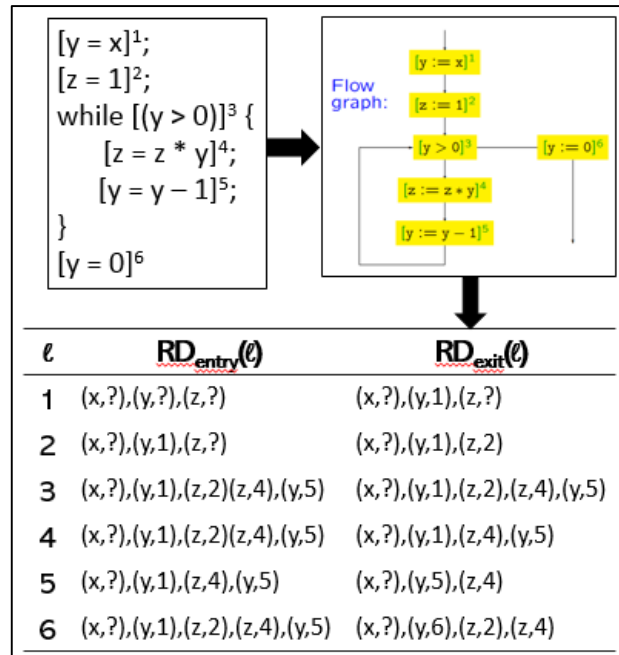
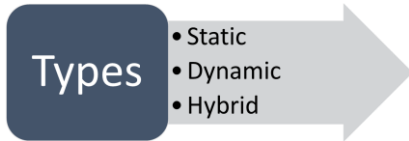
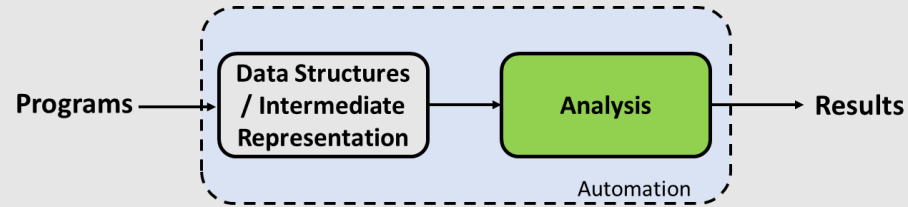
venkateshv@cmi.ac.in
Mar – Apr, 2018
Chennai Mathematical Institute



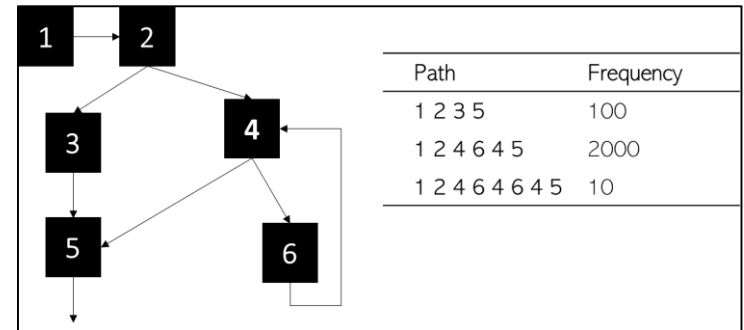
“The supply of grand challenges ... shows little sign of drying up.”

– Harman and O’Hearn in “Opportunities and Open Problems for Static and Dynamic Program Analysis”, Madrid, Spain, 2018.

Quick Review



Static Analysis

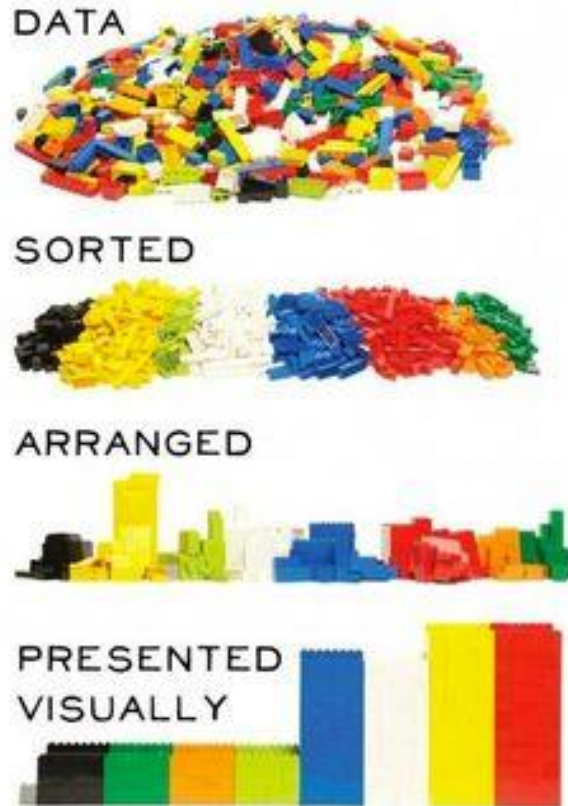


Dynamic Analysis

Program Representations

Partially covered in *Sections 1.1, 1.2, 2.1, 6.2.1 and 6.2.4 of Dragon Book.*

Finding Order in Chaos



Source: <https://visual.ly/community/infographic/how/what-infographic> by Hot Butter Studio.

Agenda

- Anatomy of a Program
 - Lexemes
 - Tokens
 - Abstract Syntax Trees (AST)
 - Intermediate Representations
- AST and CFG
 - Hands-On Session
- Popular Program Representations

Identify a Java Program

1

```
public class Fibonacci {  
  
    public static void main(String[] args) {  
  
        int n = 10, t1 = 0, t2 = 1;  
        System.out.print("First " + n + " terms: ");  
  
        for (int i = 1; i <= n; ++i)  
        {  
            System.out.print(t1 + " ");  
  
            int sum = t1 + t2;  
            t1 = t2;  
            t2 = sum;  
        }  
    }  
}
```

2

java.lang

Class Object

java.lang.Object

public class Object

Class Object is the root of the class hierarchy. Every class has Object as a superclass. All objects, including arrays, implement the methods of this class.

Since:

JDK1.0

See Also:

Class

3



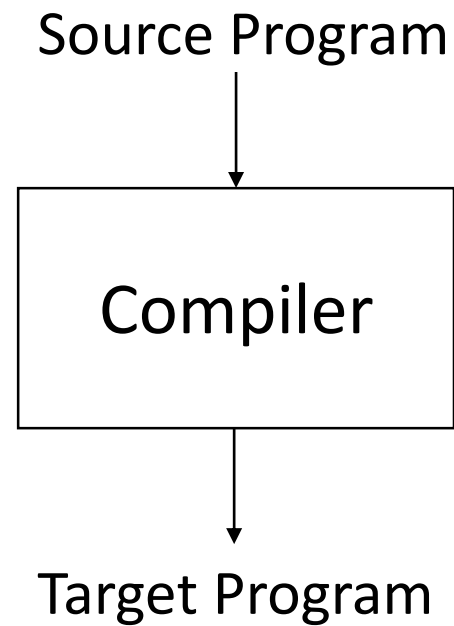
A Simple Grammar

Our programs are usually defined by a simple grammar...

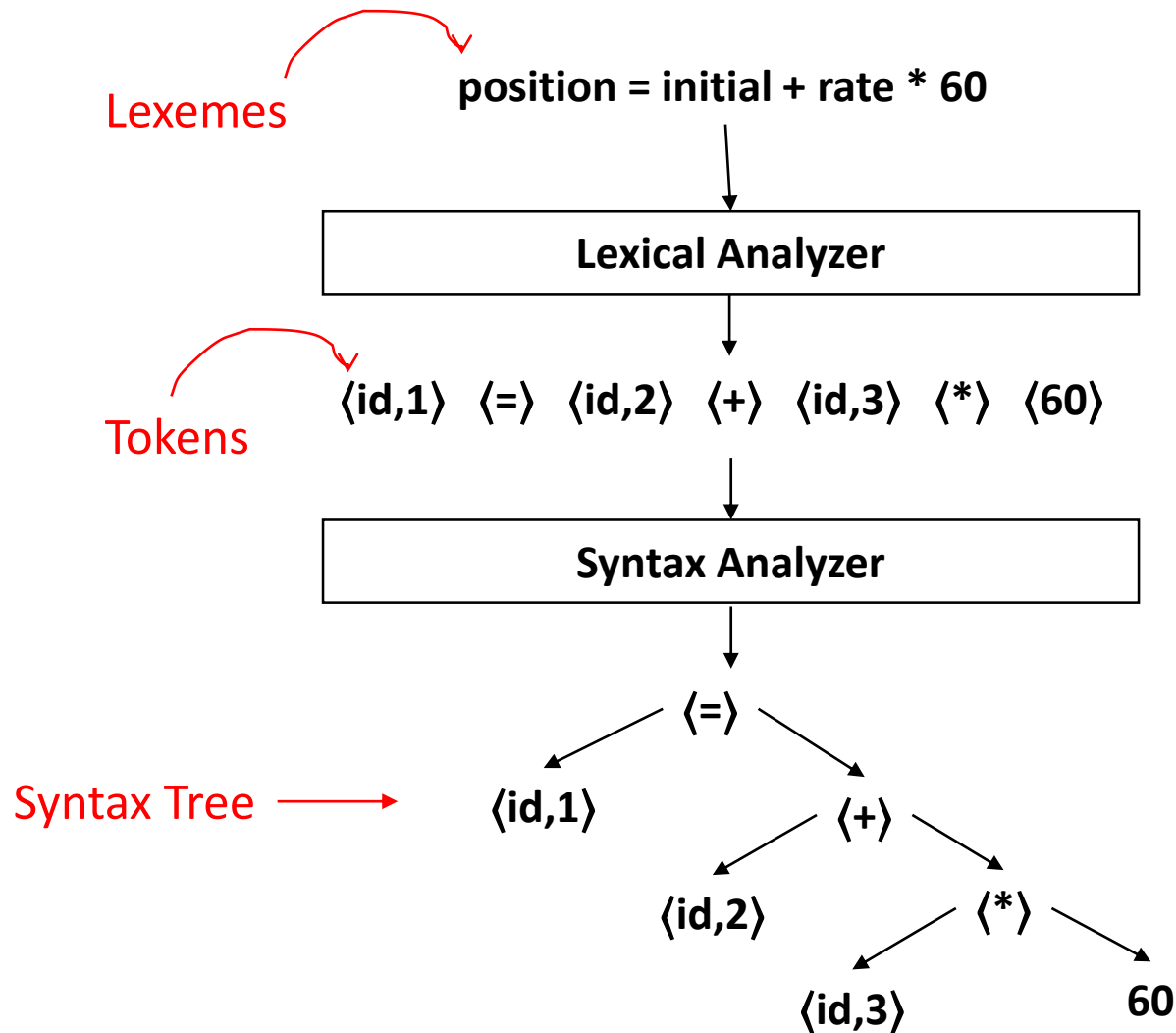
```
a ::= x | n | a1 opa a2
b ::= true | false | not b | b1 opb b2 | a1 opr a2
S ::= x := a | skip | S1; S2 |
    if (b) then S1 else S2 | while (b) do S
```

S refers to Statements, **a** is an Arithmetic Expression, and **b** is a Boolean Expression

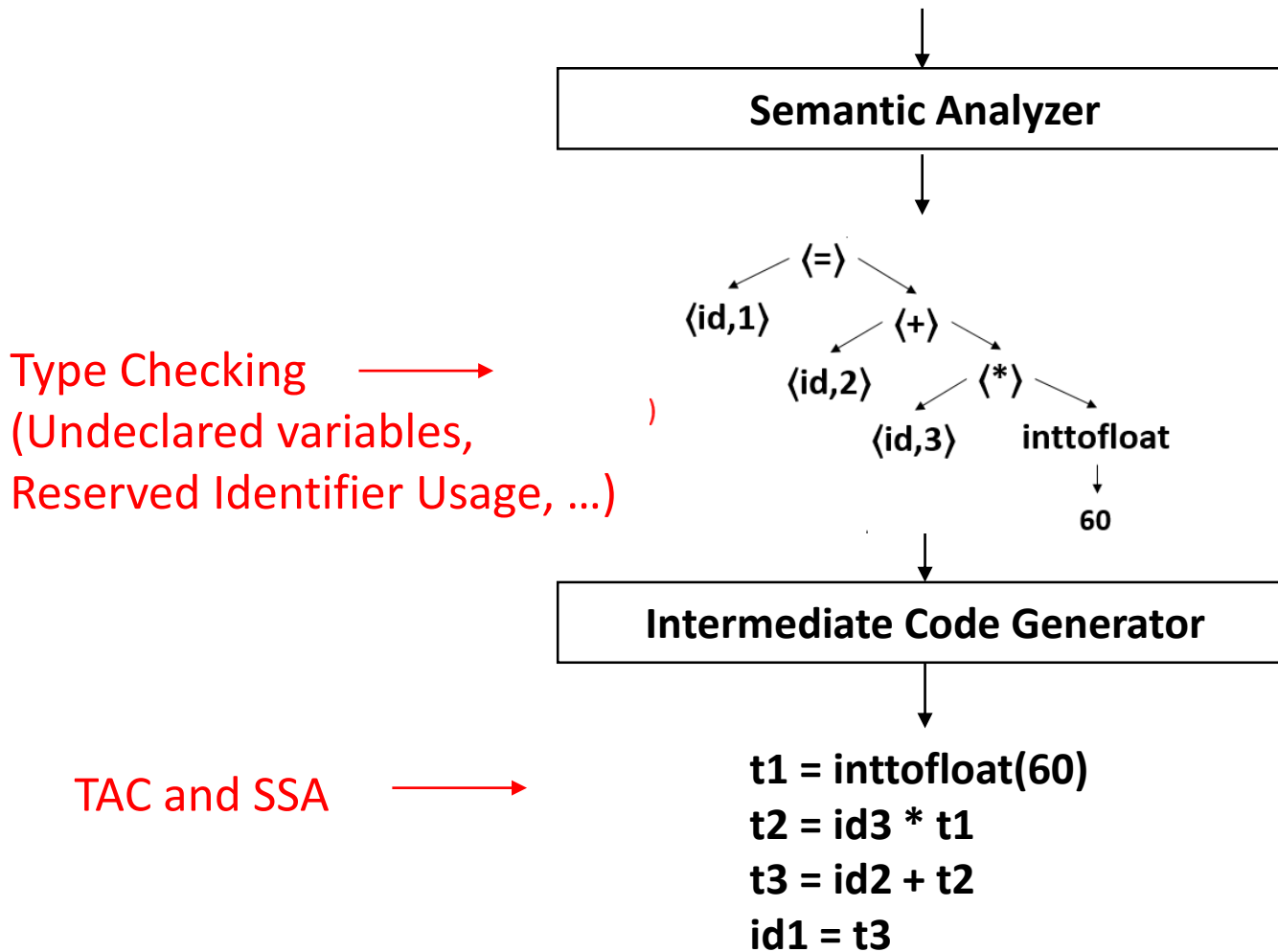
Compiler



Structure of a Compiler

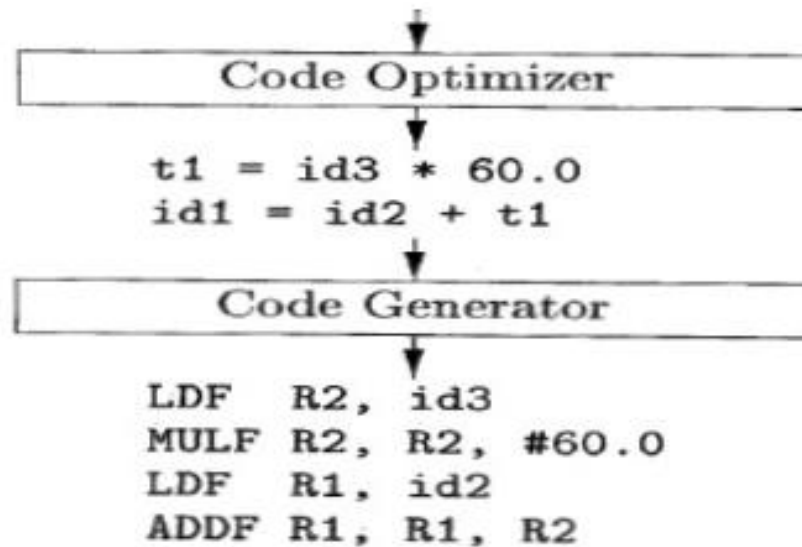


Structure of a Compiler...



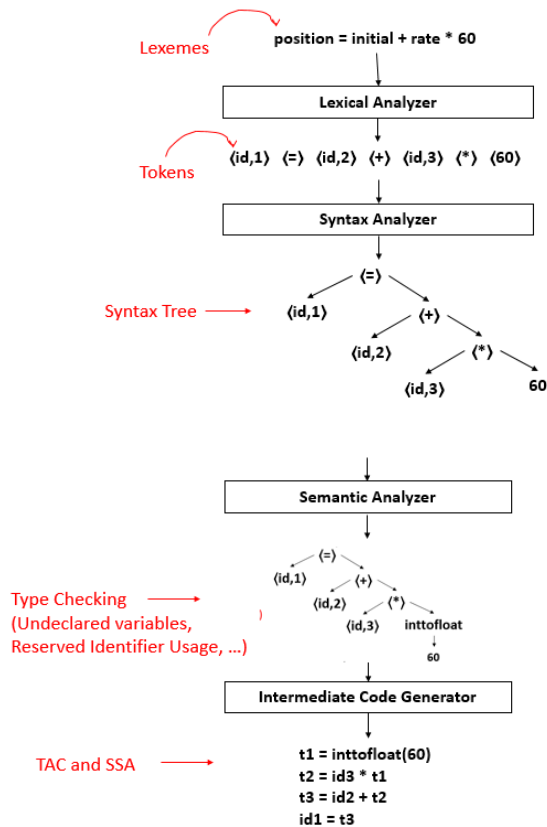
Structure of a Compiler

Program
Analysis →



“Structure of a Compiler”, taken from the dragon book.

Phases



Present day compilers take multiple passes
They enrich the IR in each pass.

Our Interests

1. Can we leverage the compiler infrastructure to analyze programs?
2. Can we improve compiler optimizations?

What is the Output?

```
public class Test {  
    public static void main(String[] args)  
    {  
        String str = "CMI";  
        StringBuilder sb = new StringBuilder();  
        for(int i = str.length() - 1; i >= 0; i--)  
        {  
            if (i == 0) sb.append('S');  
            sb.append(str.charAt(i));  
        }  
        System.out.println(sb.toString());  
    }  
}
```

Graded Quiz: Parsing

- Is this grammar ambiguous? Why or Why not?

$$S \rightarrow AS \mid \epsilon$$
$$A \rightarrow A1 \mid 0A1 \mid 01$$


Our First Analysis

using Eclipse JDT

Source Code as a Tree

- Abstract Syntax Tree

Code Snippet

```
while b  $\neq$  0  
if a > b  
  a := a - b  
else  
  b := b - a  
return a
```

AST

- **while**
 - Op: \neq
 - Var: **b**
 - Constant: **0**
 - Body
 - Branch
 - ...
- Return
 - Var: **a**

AST for Java Code

- Install Eclipse
- Install the Eclipse plugin AST View (from <https://www.eclipse.org/jdt/ui/astview/index.php>)
- Write any Java Code
- Follow the instruction in the AST View web page to view the AST.

An Example

```
1 public class VariableValueAnalysis {
2     public static void main(String[] args) {
3         int x = 10;
4         while(x > 5) {
5             x--;
6         }
7         System.out.println(x);
8     }
9 }
```

VariableValueAnalysis.java (AST Level 8). Creation time: 30 ms. Size: 36 nodes, 4,356 bytes (/

